

# New Methods for Computing Roots of Polynomials

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## 1. DIFFICULTIES OF COMPUTING POLYNOMIAL ROOTS

There exist many algorithms for computing the roots of a polynomial:

- Bairstow, Graeffe, Jenkins-Traub, Laguerre, Müller, Newton, . . .

These methods yield satisfactory results if:

- The polynomial is of moderate degree
- The roots are simple and well-separated
- A good starting point in the iterative scheme is used

This heuristic has exceptions:

$$f(x) = \prod_{i=1}^{20} (x - i) = (x - 1)(x - 2) \cdots (x - 20)$$

**Example 1.1** Consider the polynomial

$$x^4 - 4x^3 + 6x^2 - 4x + 1 = (x - 1)^4$$

whose root is  $x = 1$  with multiplicity 4. MATLAB returns the roots

1.0002, 1.0000 + 0.0002i, 1.0000 - 0.0002i, 0.9998



**Example 1.2** The roots of the polynomial  $(x - 1)^{100}$  were computed by MATLAB.

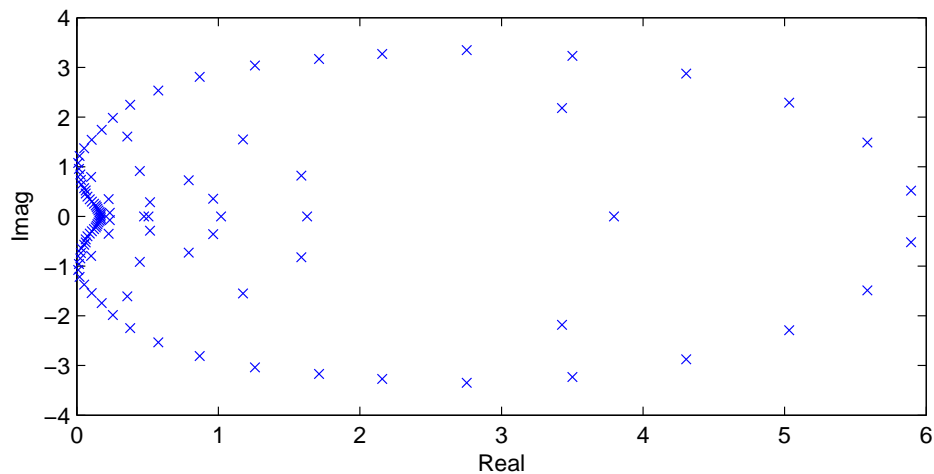


Figure 2.1: The computed roots of  $(x - 1)^{100}$ .



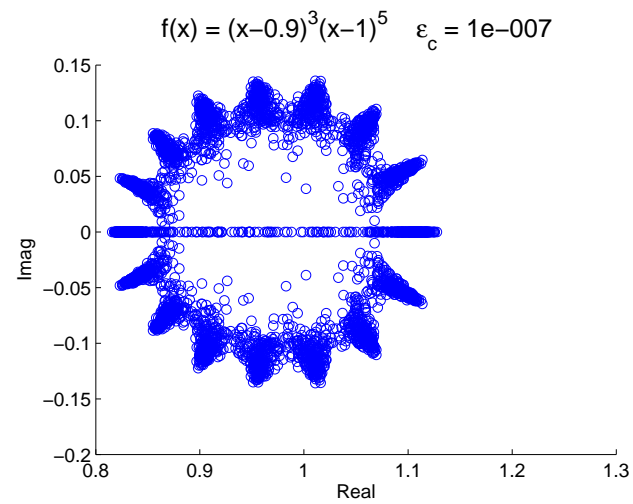
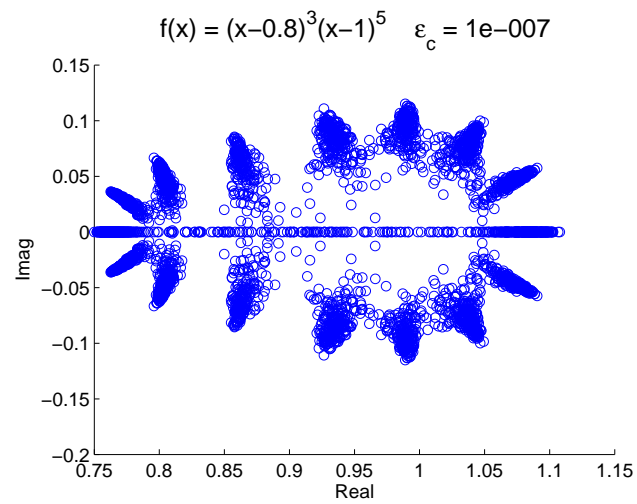
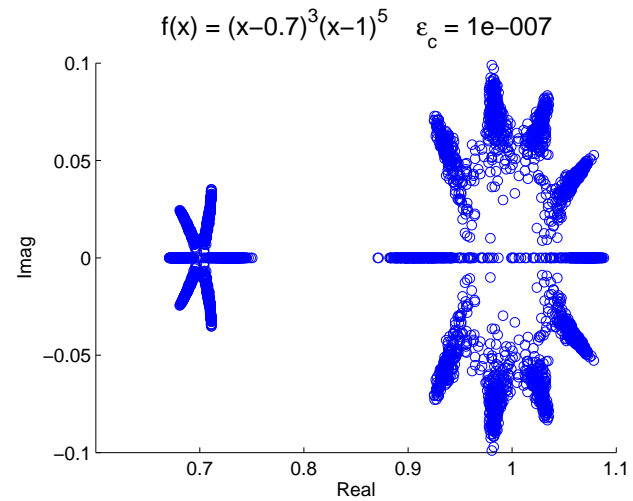
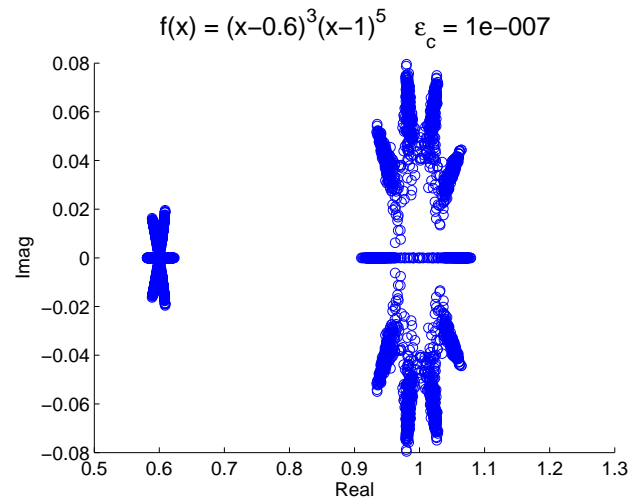


Figure 1.1: The root distribution of four perturbed polynomials.

## 2. THE GEOMETRY OF ILL-CONDITIONED POLYNOMIALS

- A root  $x_0$  of multiplicity  $r$  introduces  $(r - 1)$  constraints on the coefficients.
- A monic polynomial of degree  $m$  has  $m$  degrees of freedom.
- The root  $x_0$  lies on a manifold of dimension  $(m - r + 1)$  in a space of dimension  $m$ .
- This manifold is called a **pejorative manifold** because polynomials near this manifold are ill-conditioned.
- A polynomial that lies on a pejorative manifold is well-conditioned with respect to (the structured) perturbations that keep it on the manifold, which corresponds to the situation in which the multiplicity of the roots is preserved.
- A polynomial is ill-conditioned with respect to perturbations that move it off the manifold, which corresponds to the situation in which a multiple root breaks up into a cluster of simple roots.

Example 2.1 Consider a cubic polynomial  $f(x)$  with real roots  $x_0, x_1$  and  $x_2$

$$(x - x_0)(x - x_1)(x - x_2) = x^3 - (x_0 + x_1 + x_2)x^2 + (x_0x_1 + x_1x_2 + x_2x_0)x - x_0x_1x_2$$

- If  $f(x)$  has one double root and one simple root, then  $x_0 = x_1 \neq x_2$  and thus  $f(x)$  can be written as

$$x^3 - (2x_1 + x_2)x^2 + (x_1^2 + 2x_1x_2)x - x_1^2x_2$$

The pejorative manifold of a cubic polynomial that has a double root is the surface defined by

$$\left( \begin{array}{ccc} -(2x_1 + x_2) & (x_1^2 + 2x_1x_2) & -x_1^2x_2 \end{array} \right) \quad x_1 \neq x_2$$

- If  $f(x)$  has a triple root, then  $x_0 = x_1 = x_2$  and thus  $f(x)$  can be written as

$$x^3 - 3x_0x^2 + 3x_0^2x - x_0^3$$

The pejorative manifold of a cubic polynomial that has a triple root is the curve defined by

$$\left( \begin{array}{ccc} -3x_0 & 3x_0^2 & -x_0^3 \end{array} \right)$$





**Theorem 2.1** The condition number of the real root  $x_0$  of multiplicity  $r$  of the polynomial  $f(x) = (x - x_0)^r$ , such that the perturbed polynomial also has a root of multiplicity  $r$ , is

$$\rho(x_0) := \frac{\Delta x_0}{\Delta f} = \frac{1}{r |x_0|} \frac{\|(x - x_0)^r\|}{\|(x - x_0)^{r-1}\|} = \frac{1}{r |x_0|} \left( \frac{\sum_{i=0}^r \binom{r}{i}^2 (x_0)^{2i}}{\sum_{i=0}^{r-1} \binom{r-1}{i}^2 (x_0)^{2i}} \right)^{\frac{1}{2}}$$

where  $\|\cdot\| = \|\cdot\|_2$  and

$$\Delta f = \frac{\|\delta f\|}{\|f\|} \quad \text{and} \quad \Delta x_0 = \frac{|\delta x_0|}{|x_0|}$$

Example 2.2 The condition number  $\rho(1)$  of the root  $x_0 = 1$  of  $(x - 1)^r$  is

$$\rho(1) = \frac{1}{r} \left( \frac{\sum_{i=0}^r \binom{r}{i}^2}{\sum_{i=0}^{r-1} \binom{r-1}{i}^2} \right)^{\frac{1}{2}}$$

This expression reduces to

$$\rho(1) = \frac{1}{r} \sqrt{\frac{\binom{2r}{r}}{\binom{2(r-1)}{r-1}}} = \frac{1}{r} \sqrt{\frac{2(2r-1)}{r}} \approx \frac{2}{r} \quad \text{for large } r$$

Compare with the componentwise and normwise condition numbers

$$\kappa_c(1) \approx \frac{|\delta x_0|}{\varepsilon_c} \quad \text{and} \quad \kappa_n(1) \approx \frac{|\delta x_0|}{\varepsilon_n}$$

- $\rho(1)$  is independent of the the noise level (assumed to be small)
- $\rho(1)$  decreases as the multiplicity  $r$  of the root  $x_0 = 1$  increases



### 3. A SIMPLE POLYNOMIAL ROOT FINDER

Let  $w_i(x)$  be the product of all factors of degree  $i$  of  $f(x)$

$$f(x) = w_1(x)w_2^2(x)w_3^3(x) \cdots w_{r_{\max}}^{r_{\max}}(x)$$

Perform a sequence of GCD computations

$$\begin{aligned} q_1(x) &= \text{GCD} \left( f(x), f^{(1)}(x) \right) = w_2(x)w_3^2(x)w_4^3(x) \cdots w_{r_{\max}}^{r_{\max}-1}(x) \\ q_2(x) &= \text{GCD} \left( q_1(x), q_1^{(1)}(x) \right) = w_3(x)w_4^2(x)w_5^3(x) \cdots w_{r_{\max}}^{r_{\max}-2}(x) \\ q_3(x) &= \text{GCD} \left( q_2(x), q_2^{(1)}(x) \right) = w_4(x)w_5^2(x)w_6^3(x) \cdots w_{r_{\max}}^{r_{\max}-3}(x) \\ q_4(x) &= \text{GCD} \left( q_3(x), q_3^{(1)}(x) \right) = w_5(x)w_6^2(x)w_7^3(x) \cdots w_{r_{\max}}^{r_{\max}-4}(x) \\ &\vdots \end{aligned}$$

The sequence terminates at  $q_{r_{\max}}(x)$ , which is a constant.

A set of polynomials  $h_i(x)$ ,  $i = 1, \dots, r_{\max}$ , is defined such that

$$h_1(x) = \frac{f(x)}{q_1(x)} = w_1(x)w_2(x)w_3(x) \cdots$$

$$h_2(x) = \frac{q_1(x)}{q_2(x)} = w_2(x)w_3(x) \cdots$$

$$h_3(x) = \frac{q_2(x)}{q_3(x)} = w_3(x) \cdots$$

$\vdots$

$$h_{r_{\max}}(x) = \frac{q_{r_{\max}-2}}{q_{r_{\max}-1}} = w_{r_{\max}}(x)$$

The functions,  $w_1(x)$ ,  $w_2(x)$ ,  $\cdots$ ,  $w_{r_{\max}}(x)$ , are determined from

$$w_1(x) = \frac{h_1(x)}{h_2(x)}, \quad w_2(x) = \frac{h_2(x)}{h_3(x)}, \quad \cdots, \quad w_{r_{\max}-1}(x) = \frac{h_{r_{\max}-1}(x)}{h_{r_{\max}}(x)}$$

until

$$w_{r_{\max}}(x) = h_{r_{\max}}(x)$$

The equations

$$w_1(x) = 0, \quad w_2(x) = 0, \quad \dots, \quad w_{r_{\max}}(x) = 0$$

contain only simple roots, and they yield the simple, double, triple, etc., roots of  $f(x)$ .

- If  $x_0$  is a root of  $w_i(x)$ , then it is a root of multiplicity  $i$  of  $f(x)$ .



Mathematical operations performed in this root finder:

- GCD computations
- Polynomial division
- Solution of simple polynomial equations

### 3.1 Discussion of method

- The computation of the GCD of two polynomials is an ill-posed problem because it is not a continuous function of their coefficients:
  - The polynomials  $f(x)$  and  $g(x)$  may have a non-constant GCD, but the perturbed polynomials  $f(x) + \delta f(x)$  and  $g(x) + \delta g(x)$  may be coprime.
- The determination of the degree of the GCD of two polynomials reduces to the determination of the rank of a resultant matrix, but the rank of a matrix is not defined in a floating point environment. The determination of the rank of a noisy matrix is a challenging problem that arises in many applications.
- Polynomial division, which reduces to the deconvolution of their coefficients, is an ill-conditioned problem that must be implemented with care in order to obtain a computationally reliable solution.

#### 4. APPROXIMATE GREATEST COMMON DIVISORS

If  $f(x)$  is exact and all computations are performed in a symbolic environment, the GCD of  $f(x)$  and its derivative  $f^{(1)}(x)$  can be computed by the Sylvester resultant matrix  $S(f, f^{(1)})$ .

The polynomial  $f(x)$  is rarely known exactly, and so the given data is

$$\tilde{f}(x) = f(x) + \delta f(x)$$

and  $\tilde{f}(x)$  and  $\tilde{f}^{(1)}(x)$  are (with probability almost 1) coprime.

- The polynomials  $\tilde{f}(x)$  and  $\tilde{f}^{(1)}(x)$  have an **approximate greatest common divisor**.
- Use the method of **structured total least norm** applied to  $S(\tilde{f}, \tilde{f}^{(1)})$  to compute the smallest perturbation of  $S(\tilde{f}, \tilde{f}^{(1)})$  such that its perturbed form is singular, which implies that the perturbed form of  $\tilde{f}(x)$  has a multiple root.





#### 4.1 The non-uniqueness of the Sylvester resultant matrix

- An approximate GCD of  $f(x)$  and  $g(x)$  is equal to, up to a scalar multiplier, an approximate GCD of  $f(x)$  and  $\alpha g(x)$ , where  $\alpha$  is an arbitrary non-zero constant.
- The resultant matrix  $S(f, \alpha g)$  should be used when it is desired to compute an approximate GCD of  $f(x)$  and  $g(x)$ .
- Since  $S(f, \alpha g) \neq \alpha S(f, g)$ , the inclusion of  $\alpha$  permits a family of approximate GCDs, rather than only one approximate GCD, to be computed.

Example 4.2 Consider the exact polynomials

$$\hat{f}_1(x) = (x - 0.25)^8(x - 0.5)^9(x - 0.75)^{10}(x - 1)^{11}(x - 1.25)^{12}$$

$$\hat{g}_1(x) = (x + 0.25)^4(x - 0.25)^5(x - 0.5)^6$$

with signal-to-noise ratio  $\mu = 10^8$ .

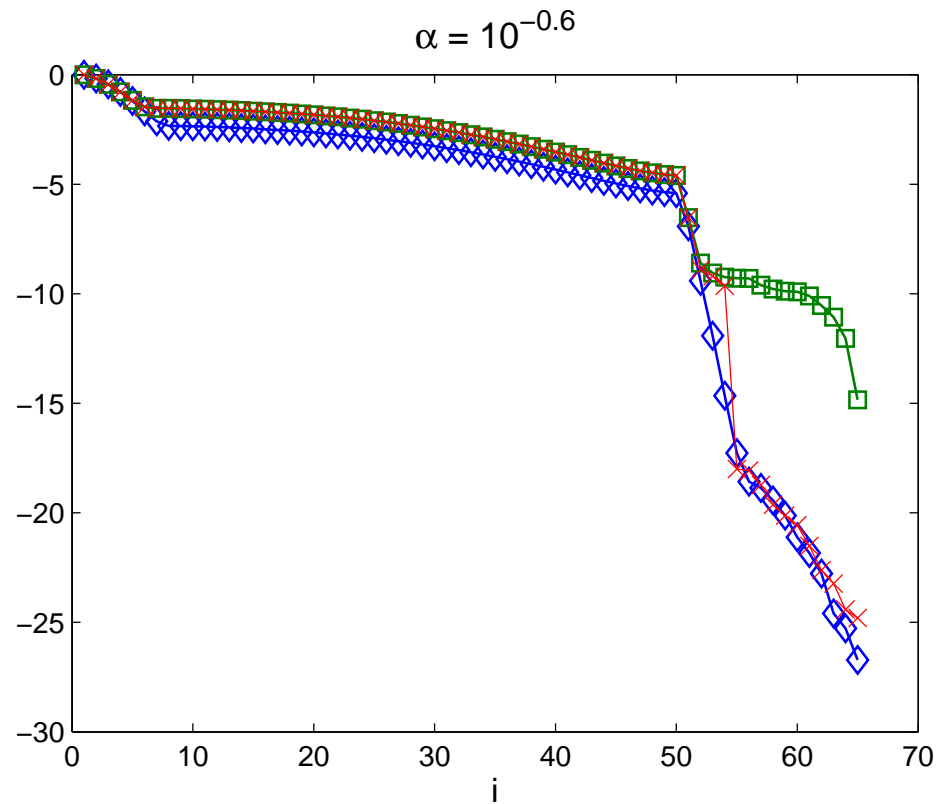


Figure 4.1: The normalised singular values of the Sylvester resultant matrix, on a logarithmic scale, for (i) the theoretically exact data  $S(\hat{f}_1, \hat{g}_1)$ ,  $\diamond$ ; (ii) the given inexact data  $S(f_1, g_1)$ ,  $\square$ ; (iii) the computed data  $S(\tilde{f}_{1,0}, \tilde{g}_{1,0})$ ,  $\times$ , for  $\alpha = 10^{-0.6}$ .  $\square$

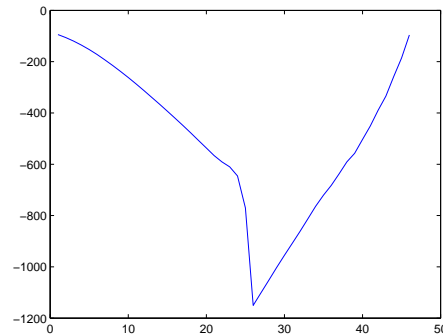
## 5. THE CALCULATION OF THE POLYNOMIAL ROOTS

- Use the method of least squares to perform the polynomial division
- Obtain initial estimates of the roots of the polynomial by solving a set of polynomial equations, each of whose roots is simple.
- Refine these estimates by using the method of non-linear least squares.

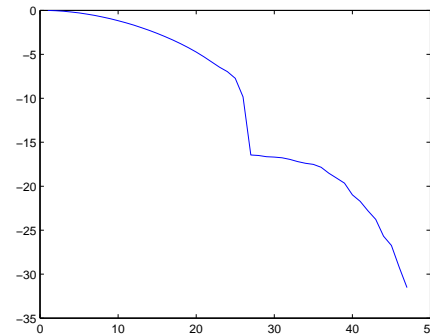
## 6. EXAMPLES

**Example 6.1** Consider the Sylvester resultant matrix of  $f(x)$  and its derivative  $f^{(1)}(x)$ , which has rank 26.

$$f(x) = (x - 0.1)^8(x - 0.5)^8(x - 0.9)^8$$

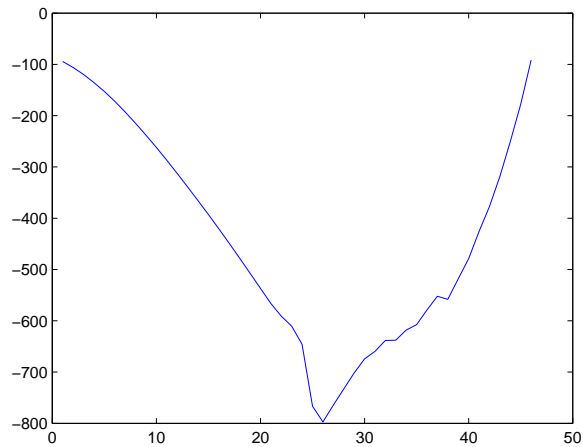


(i)

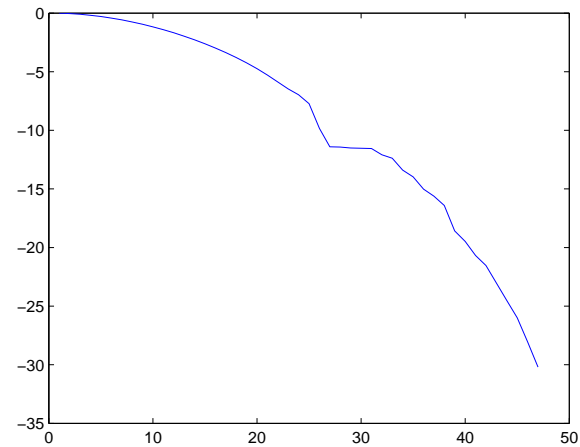


(ii)

Figure 6.1: (i) The rank estimate of  $S(f, f^{(1)})$  from the principle of maximum likelihood, and (ii) the singular values of  $S(f, f^{(1)})$ , in the absence of noise.



(i)



(ii)

Figure 6.2: (i) The rank estimate of  $S(f, f^{(1)})$  from the principle of maximum likelihood, and (ii) the singular values of  $S(f, f^{(1)})$ , for a signal-to-noise ratio of  $10^9$ .



Example 6.2 Consider the polynomial

$$f(x) = (x - 1)^{20}(x - 2)^{15}(x - 3)^{10}(x - 4)^5$$

The computed roots are

Multiplicity	Root
1	20
2	15
3	10
4	5

## 7. FURTHER RESEARCH

- Use the displacement structure of matrices to optimise the computational efficiency of the method.
- Optimise the calculation of the scale parameter  $\alpha$  in the Sylvester resultant matrix.
- Investigate the use of threshold independent methods, for example, the principle of maximum likelihood, for the estimation of the rank of a matrix.
- Consider the problem that occurs when there are bounds on the displacement of each coefficient.
- Extend to bivariate and trivariate polynomials.



## 8. SUMMARY

- A radically new method of solving polynomial equations has been described.
- This method first calculates the multiplicities of the roots, and then computes their values.
- The method is a computational implementation of the theory of pejorative manifolds.
- Computational experiments show that it is very successful.