

## **Self-dual spin 1/2 systems, zeros of partition function, error correcting codes**

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The last three pages of the present file have figures (‘end-figures’) with marked zeros of certain polynomials: partition functions of two (strongly) self-dual spin 1/2 systems with periodic boundary conditions and (Hamming) weight enumerators of  $([14, 7, 5]$  and  $[30, 5, 6])$  cyclic self-dual binary (error correcting) codes.

Apart from the case of  $TT[6, 6]$ , all the zeros, with the exception of a few, are located on two circles (‘F-circles’): in the case of  $A6[6, 10]$  only one zero (modulo the symmetry group of the figure) is off the F-circles.

In the following, we will discuss the systems in question, and we will argue that a unified treatment of both of them yields new insights into the problem of the location of the zeros and new (and interesting) codes.

## Spin 1/2 Systems

The *energy* of a (*transitive, translation invariant*) spin 1/2 ferromagnetic system on the lattice  $\mathbb{Z}^d$  has the form

$$H = - \sum_{B \in \mathcal{B}} K(B) \sigma_B, \quad K(B) > 0,$$

where

$$\sigma_B = \prod_{a \in B} \sigma_a, \quad \sigma_a = \pm 1,$$

$\mathcal{B}$  is a translation invariant family of finite subsets, with a finite fundamental subfamily  $\mathcal{B}_0$  generating it by translations. The *coupling constants (interaction)*  $K(B)$  are proportional to inverse temperature.

Univariate version: all the coupling constants are equal.

**Examples:** The figures show a *fundamental family*  $\mathcal{B}_0$  of *bonds*.

– Ising Model (IM)



– TT Model (solved by Baxter and Wu)



– A6 Model



• **Partition function.** For a finite subsystem  $\Lambda$  of the infinite system,

$$\begin{aligned} Z &= Z_\Lambda(K) = \sum_{\sigma} \exp(-H_\Lambda(\sigma)) \\ &= \sum_{\sigma_a = \pm 1, a \in \Lambda} \exp \sum_{B \in \mathcal{B}_\Lambda} K(B) \sigma_B, \end{aligned}$$

$\mathcal{B}_\Lambda := \{B \in \mathcal{B} : B \subset \Lambda\}$ . Here finite volumes  $\Lambda$  will be taken with *periodic boundary conditions* (‘periodic version of the model’). For example, we may take  $\Lambda = \mathbb{Z}_m \times \mathbb{Z}_n$  in two dimensions. We will denote this periodic version by suffix  $[m, n]$ : TT[6, 9] is the TT model on  $\Lambda = \mathbb{Z}_6 \times \mathbb{Z}_9$ , with periodic boundary conditions.

■ **The main result for IM** (1941). Define the *dual* interaction  $\mathbf{K}^* = (K_1^*, K_2^*)$  by

$$\tanh K_1^* = \exp -2K_2, \quad \tanh K_2^* = \exp -2K_1$$

Then

$$P(\mathbf{K}) - \frac{1}{2} \ln \prod_i \cosh 2K_i = P(\mathbf{K}^*) - \frac{1}{2} \ln \prod_i \cosh 2K_i^*$$

$$(P = \lim_{\Lambda} |\Lambda|^{-1} \ln Z_{\Lambda})$$

**Corollary** (location of the critical point). *For equal coupling constants,  $K_1 = K_2 = K$*

$$\tanh K = \exp -2K \quad \Rightarrow \quad K_{\text{cr}} = \frac{1}{2} \ln (\sqrt{2} + 1)$$

(This is the IM first exact result, showing, in particular, that a number of results obtained through approximate calculations cannot be true)

## General Formulation of Duality (1971+).

For any finite set  $\Lambda$  and a family  $\mathcal{B}$  of its subsets, introduce

configuration space  $\Omega = \Omega_\Lambda := \{-1, 1\}^\Lambda$ ;

$\gamma : \Omega \rightarrow \mathcal{P}(\mathcal{B})$  (contour map),

$$\gamma(\sigma) := \{B \in \mathcal{B} : \sigma_B = -1\}, \quad \Gamma := \text{Im}(\gamma)$$

$$\mathcal{K} := \{\beta \subset \mathcal{B} : \prod_{B \in \beta} \sigma_B = 1, \forall \sigma \in \Omega\}$$

$$\mathcal{S} := \{\sigma : \sigma_B = 1, \forall B \in \mathcal{B}\} = \{\sigma : \gamma(\sigma) = \emptyset\}$$

elements of  $\Gamma$ ,  $\mathcal{K}$ ,  $\Omega$  and  $\mathcal{S}$  are called *contours*, *cycles*, *configurations* and *ground states*, respectively.  $|\mathcal{S}| = 2$  for IM,  $|\mathcal{S}| = 4$  for the TT model, and  $|\mathcal{S}| = 16$  for the A6 model. These definitions are extended to infinite  $\Lambda$  in appropriate way.

- **Low and High Temperature Expansions** Introduce the polynomials

$$Z^{\text{LT}} = Z_{\Lambda}^{\text{LT}}(w) = \sum_{\beta \in \Gamma} w^{\beta}, \quad w^{\beta} = \prod_{B \in \beta} w_B,$$

and

$$Z^{\text{HT}} = Z_{\Lambda}^{\text{HT}}(x) = \sum_{\beta \in \mathcal{K}} x^{\beta}, \quad x^{\beta} = \prod_{B \in \beta} x_B.$$

When  $w_B = e^{-2K(B)}$  and  $x_B = \tanh K(B)$  both are equal to the partition function, up to trivial factors, yielding, respectively, its Low and High Temperature Expansions.

(Finite) system is *self-dual* when the two polynomials are equal, *strongly self-dual* if  $\mathcal{K} = \Gamma$



With the natural group structures on  $\Omega$  and  $\mathcal{P}(\mathcal{B}) \simeq \mathbb{F}_2^{\mathcal{B}}$  ( $\mathbb{F}_2 = \{0, 1\}$  is the two-element field),  $\sigma_B$  become characters;  $\gamma$  a homomorphism;  $\mathcal{S} \subset \Omega$ ,  $\Gamma \subset \mathcal{P}(\mathcal{B})$  and  $\mathcal{K} \subset \mathcal{P}(\mathcal{B})$  subgroups.

Moreover, for  $\Lambda = \mathbb{Z}^d$ , translations act in a natural way by homomorphisms on each of these groups. For infinite systems one has:

**Definition.** A system is *self-dual* if there is a bijection  $\mathcal{B} \rightarrow \mathcal{B}$  which *commutes* or *anticommutes* with translations and defines an isomorphism of  $\Gamma$  onto  $\mathcal{K}$ . A system is *strongly self-dual* if  $\mathcal{K} = \Gamma$ .

**Theorem.** A spin 1/2 system on a lattice  $\mathbb{Z}^d$ , any dimension  $d$ , is self-dual if and only if  $|\mathcal{B}_0| = 2$ . It is strongly self-dual if and only, with  $\mathcal{B}_0 = \{B_1, B_2\}$ ,  $B_2$  is a translation of  $I(B_1)$ , where

$$I(B) = \{-a : a \in B\}.$$

Thus, while all three models of p. 4 are self-dual, only the last two are strongly self-dual.

## F(isher's)-zeros

These are the zeros of  $Z^{\text{LT}}(w)$  – the main object of the present study. They are, in a sense, of a more combinatorial nature than Lee-Yang zeros.

- Fisher (1964): for IM (no external field, in infinite volume limit), *F-zeros are on the circles*

$$|w - 1| = \sqrt{2}, \quad |w + 1| = \sqrt{2}$$

(F-circles). It was shown later that this holds for finite IM with suitable boundary conditions.

For finite IM with periodic boundary conditions, most of the F-zeros are off the F-circles, as in the case of TT [6,6].

- For the models considered below, the set of zeros is symmetric under the group ( $D_{16}$ ) generated by

$$Dx = \frac{1-x}{1+x}, \quad Ix = \frac{1}{x}, \quad Cx = x^*$$

Invariance under  $D$  follows from self-duality; invariance under  $I$  follows from *palindromicity* of  $Z$ .

## Strongly Self-dual Systems with periodic boundary conditions

One has  $\mathcal{B} = \mathcal{B}_\Lambda$ ,  $\Gamma, \mathcal{K} \subset \mathcal{P}(\mathcal{B}) \simeq \mathbb{F}_2^{\mathcal{B}}$ ,

$$\mathcal{K} = \Gamma^\perp, \quad \Gamma = \mathcal{K}^\perp$$

where orthogonality is with respect to the bilinear form on  $\mathbb{F}_2^{\mathcal{B}}$ :

$$(\alpha, \beta) \mapsto \alpha \cdot \beta := \sum_{B \in \mathcal{B}} \alpha_B \beta_B$$

(sum in  $\mathbb{F}_2$ ). By linear algebra,

$$\dim_{\mathbb{F}_2} \Gamma + \dim_{\mathbb{F}_2} \mathcal{K} = |\mathcal{B}| .$$

Since  $2^{|\Lambda|} = |\mathcal{S}| \cdot |\Gamma|$  ( $\mathcal{S} = \text{Ker } \gamma$ ,  $\Gamma = \text{Im } \gamma!$ ), this implies that a periodic version of a strongly self-dual model is (strongly) self-dual if and only if  $\mathcal{S}$  is trivial. Equivalently, if and only if no ground state, apart from the trivial one, has the periodicity of  $\Lambda$ . Hence one has to consider only *systems with  $|B|$  odd*.<sup>1</sup>

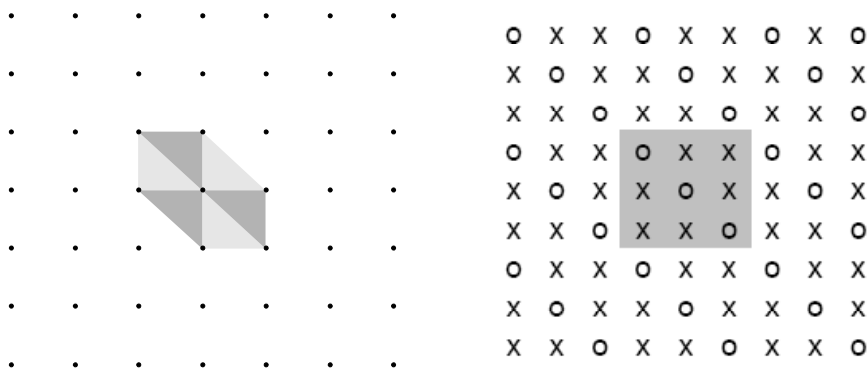
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<sup>1</sup> If  $\mathcal{S}$  is non-trivial, a quantum code that can be defined by these systems can correct  $|\mathcal{S}|^2$  errors.

In 2D, (after possibly reducing the system, in a suitable sense) the *number of ground states is finite*, and for any strongly self-dual system with  $|B|$  odd, there exist a sequence of  $\Lambda_n \rightarrow \mathbb{Z}^2$  for which the  $\Lambda_n$ -versions of the system are (strongly) self-dual.

**Example.** TT[5, 6] model is (strictly) self-dual, while TT[6, 6] or TT[6, 9] are not. More generally, TT[ $m, n$ ] model is self-dual unless both  $m$  and  $n$  are multiples of three. The end-figures show that in the TT[5, 6]-case, the zeros are on the F-circles while in the TT[6, 6]-case they are not.

The first of the figures



shows a generator of both  $\Gamma$  and  $\mathcal{K}$ , while the second one shows a ground state of the model on  $\mathbb{Z}^2$ : o stands for +1 and  $\times$  for  $-1$ . Translating it, one obtains the other two non-trivial ground states. That no TT[ $m, n$ ] system has non-trivial ground states, unless both  $m$  and  $n$  are multiples of three, follows from periodicity properties of this ground state.

## From theory of Linear Binary Codes:

*Binary code*  $C$  of length  $n =$  subspace of  $\mathbb{F}_2^n$ . For  $c \in C$ , the *weight* of  $c$  is

$\text{wt}(c) =$  number of nonzero components of  $c$ .

(*Hamming*) *weight enumerator*  $W_C(X, Y)$  of the code  $C$  is the homogeneous polynomial

$$\begin{aligned} \text{hwe}_C(X, Y) &= W_C(X, Y) = \sum_{c \in C} X^{n-\text{wt}(c)} Y^{\text{wt}(c)} \\ &= \sum_{m=0}^n N_m(C) X^{n-m} Y^m, \end{aligned}$$

where  $N_m(C)$  is the number of codewords of weight  $m$ , which is the homogeneous form of the polynomial

$$W_C(z) = \sum_{c \in C} z^{\text{wt}(c)},$$

zeros of which are shown on the end-figures.

*Dual code*

$$C^\perp = \{c \in \mathbb{F}_2^n : c \cdot c' \equiv 0, \forall c' \in C\}.$$

A code is *self-dual* if it is equal to its dual.

## Lattice associated with self-dual codes

The *lattice*  $L_C \subset \mathbb{R}^n$  is defined by integer vectors that reduce mod 2 to codewords,

$$L_C := \{2^{-1/2}v : v \in \mathbb{Z}^n, v \bmod 2 \in C\}.$$

The lattice *dual* to  $L_C$  is the lattice  $L_{C^\perp}$  of the dual code. Hence the lattice of self-dual codes are *unimodular*.

- Theta series of a code  $C$  is the theta series of  $L_C$ :

$$\Theta_C(z) = \Theta_{L_C}(z) := \sum_{k \in L_C} z^{k \cdot k}.$$

Theta series of a code is related to its weight enumerator:

$$\Theta_C = W_C(\Theta_0, \Theta_1),$$

where

$$\Theta_0(z) := \sum_{k \in \mathbb{Z}} z^{2k^2}, \quad \Theta_1(z) := \sum_{k \in \mathbb{Z}} z^{2(k+1/2)^2}.$$

For a self-dual code  $\theta_C(\tau) := \Theta_C(e^{i\pi\tau})$  is a modular form for the group  $\Gamma(2)$ .

Full weight enumerators can be expressed through Siegel modular form.

- *Cyclic codes.*  $C$  is invariant under cyclic permutations of its elements:

$$\begin{aligned} (c_0, c_1, \dots, c_{n-2}, c_{n-1}) \in C &\Rightarrow \\ (c_{n-1}, c_0, c_1, \dots, c_{n-2}) \in C & \end{aligned}$$

Equivalent formulation: For  $c \in C$ , let

$$c(T) := \sum_{i=0}^{n-1} c_i T^i \in R = \mathbb{F}_2[T] / \langle T^n - 1 \rangle$$

then  $\{c(T) : c \in C\}$  is a (principal) ideal of  $R$ . Let  $g(T)$  be a generator of the ideal; the code is self-dual if and only if

$$I(g)(T) \cdot g(T) \equiv 0 \pmod{T^n - 1}.$$

For example, the first two end-figures show zeros of  $W_{[14,7,4]}(z)$  and  $W_{[30,15,6]}(z)$ , which are the weight enumerator of the cyclic self-dual codes denoted  $[14, 7, 4]$  (length 14) and  $[30, 15, 6]$  (length 30), respectively. The ideal of the  $[14, 7, 4]$  code is generated by

$$g(T) = 1 + T + T^2 + T^3 + T^6 + T^7;$$

$I(g)(T)$ , usually denoted  $\overleftarrow{g}(T)$ , is  $1 + T + T^4 + T^5 + T^6 + T^7$ , with  $\overleftarrow{g}(T) \cdot g(T) \equiv 0 \pmod{T^{14} - 1}$ .

More examples of cyclic self-dual codes ('trivial'):  $n = 2m$ ,  $g(T) = 1 + T^n$ . AN easy calculation yields

$$\begin{aligned} W(X, Y) &= (X^2 + Y^2)^m = X^{2m}(1 + z^2)^m \\ &= X^{2m}(z - i)^m(z + i)^m = X^{2m}W(z). \end{aligned}$$

This yields  $z = \pm i$  for roots of  $W(z)$ , roots that appear on all the end-figures.



## Strongly self-dual spin systems yield (new) self-dual multi-cyclic codes

Setting  $C = \Gamma$  one obtains a binary (multi-)cyclic code of length  $|\mathcal{B}|$  ('contour codes'). The map  $\gamma : \{-1, +1\}^\Lambda \rightarrow \mathbb{F}_2^{\mathcal{B}}$  yields a coding. The code is self-dual if and only if the spin system is strongly self-dual. Translation invariance of the spin system implies multi-cyclicity of the code.

As  $|\Lambda|$  is increased, one obtains a family of self-dual regular low-density parity-check codes (LDPC), codes that are known to exhibit good performance under message-passing decoding.

The partition function of the spin system is equal to the weight enumerator of the code:

$$Z^{\text{LT}}(w) = W_\Gamma(w) = \frac{\Theta_{L_\Gamma}(w)}{\Theta_0(w)^{|\mathcal{B}|}} = \frac{\theta_{L_\Gamma}(\tau)}{\theta_0(\tau)^{|\mathcal{B}|}},$$

$w = \theta_1(\tau)/\theta_0(\tau)$ , and the celebrated MacWilliams Theorem about relation between  $W_{C^\perp}$  and  $W_C$  is an expression of LT-HT duality discussed earlier.<sup>2</sup>

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<sup>2</sup> The cyclic codes are a degenerate case of contour codes: they are obtained when  $\mathcal{B}_0$  has just one element, say  $B$ , instead of two, and the pair  $(B, \Lambda)$  satisfies the condition  $I(B) \cdot B \equiv 0 \in \mathbb{F}_2\Lambda$ .

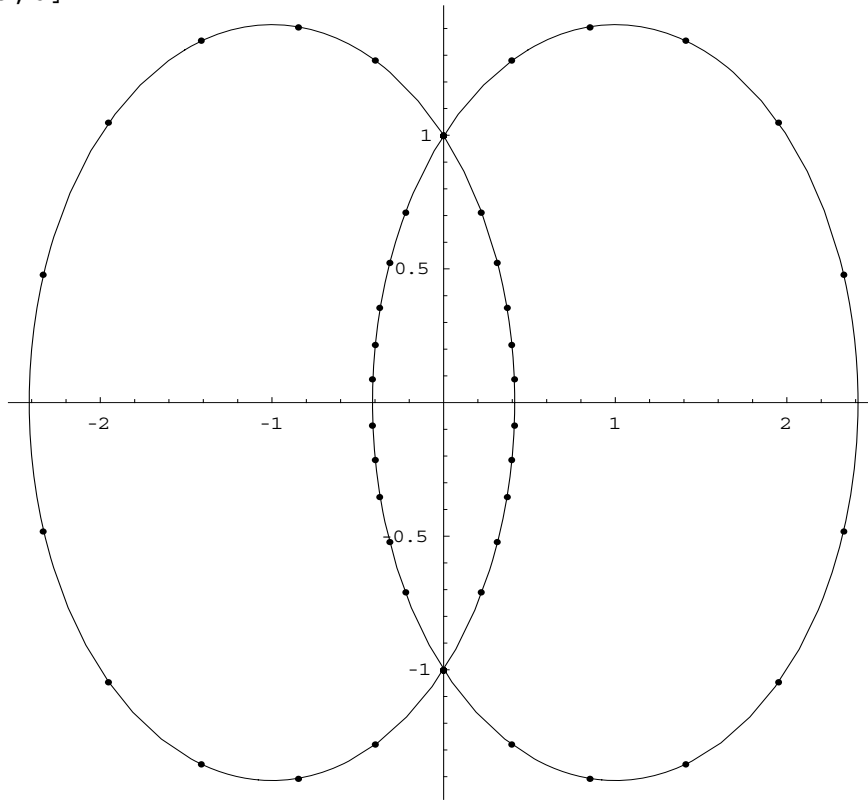
Thus, on the one hand one obtains an expression of the partition function of spin systems in terms of an exponential of a quadratic form and a modular function; on the other hand, one obtains a new family of self-dual LDPC.<sup>3</sup>

This structure (and its generalization to the non-binary case, and lattices more general than  $\mathbb{Z}^d$ ), relating spin systems and (multi-)cyclic codes, is currently being exploited in analysis of zeros of partition function and properties of the corresponding codes.

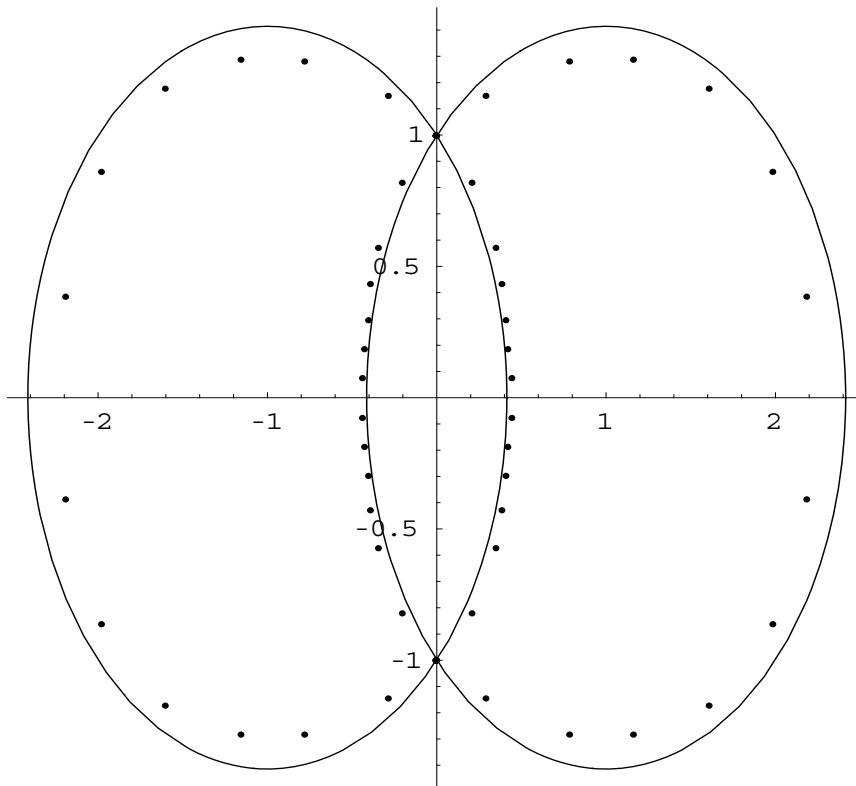
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<sup>3</sup> The counter-image of the F-circles under the map  $\tau \mapsto w$  of  $\mathbb{H}$  to  $\mathbb{C}$  consists of a nice family of semi-circles (geodesics).

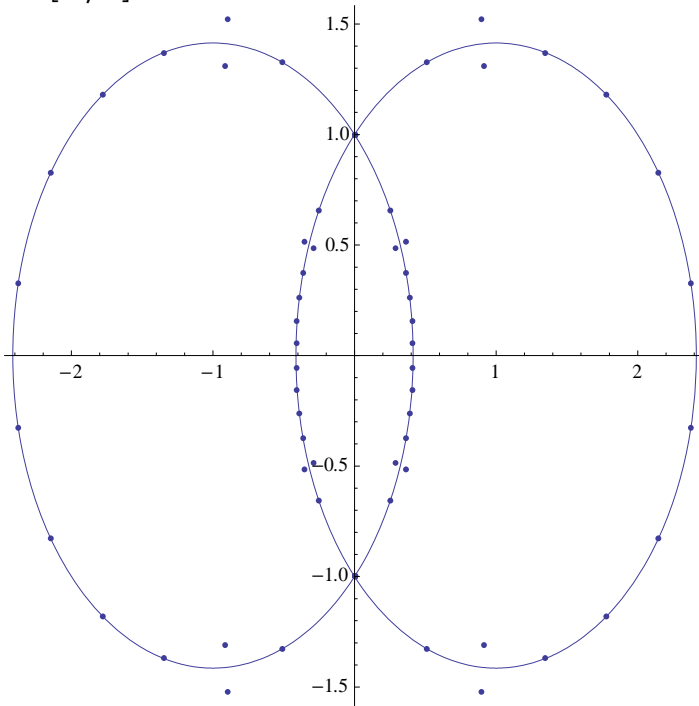
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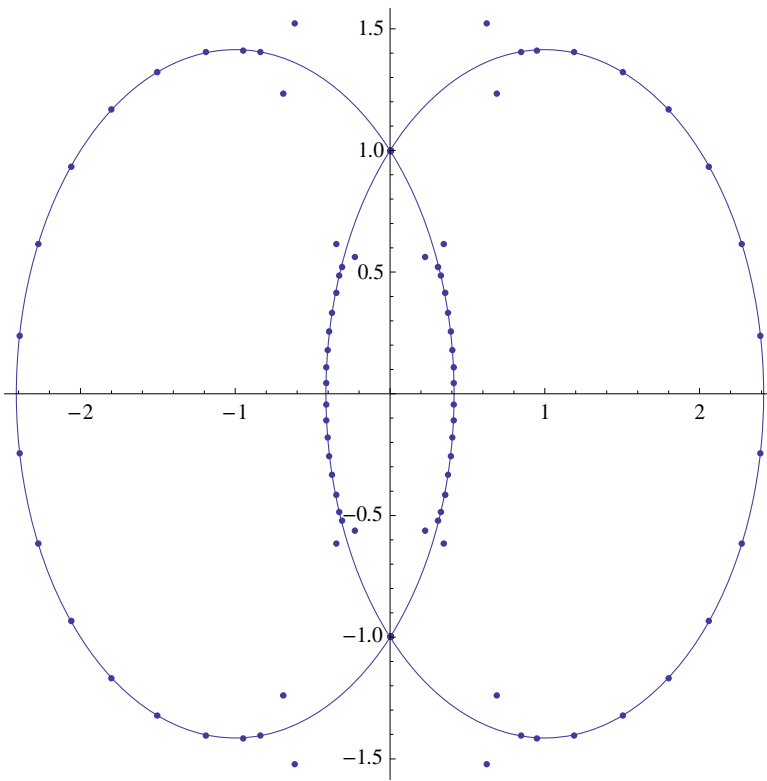
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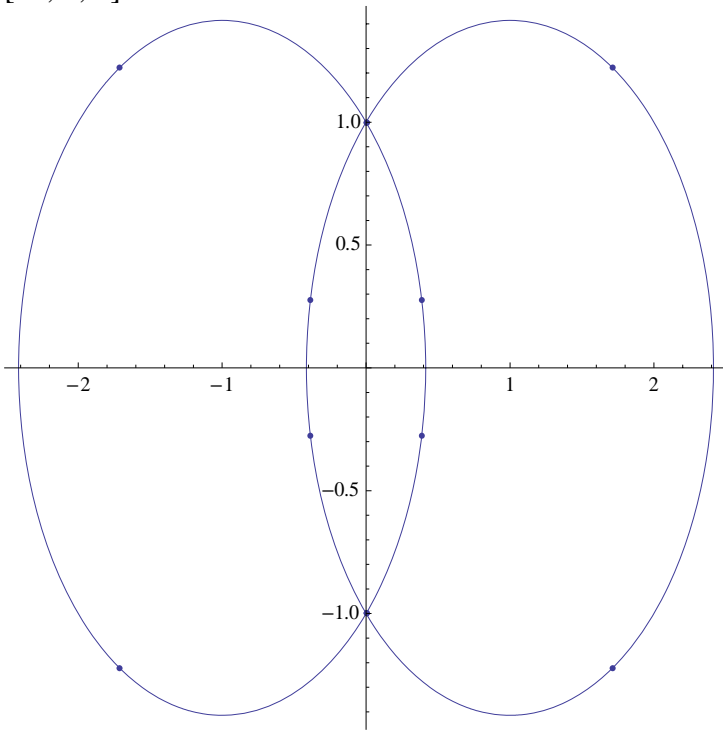
A6 [6, 7]



A6 [6, 10]



[14, 7, 4]



[30,15,6]

