

# Trees versus Connected Graphs II

William G. Faris

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# Outline

Given sum over connected graphs:

$$\sum_{G_c} \prod_{\{i,j\} \in G_c} t_{ij}.$$

Express as sum over trees

$$\sum_T \text{wt}(T).$$

How do trees emerge? From the distributive law!

# The distributive law

The product over the base of the sum over the fibers is the sum over sections of the product over the base.

$$\prod_{i \in B} \sum_{k \in F_i} a_{ik} = \sum_s \prod_{i \in B} a_{i s(i)}$$

Here  $s : B \rightarrow \bigcup_i F_i$  satisfies  $s(i) \in F_i$ .

## Example: FOIL

$$\begin{array}{cc} a & c \\ b & d \\ \downarrow & \downarrow \\ 1 & 2 \end{array}$$

$$F : 1 \mapsto a, 2 \mapsto c$$

$$O : 1 \mapsto a, 2 \mapsto d$$

$$I : 1 \mapsto b, 2 \mapsto c$$

$$L : 1 \mapsto b, 2 \mapsto d$$

$$(a + b)(c + d) = ac + ad + bc + bd$$

## Example: Interactions in physics

Potential energy between pairs of particles:

$$0 \leq v_{ij} \leq +\infty$$

Pairs of particles have probability factors:

$$0 \leq \exp(-\beta v_{ij}) = 1 + t_{ij} \leq 1.$$

Convert product to graph sum:

$$g = \prod_{\{i,j\}} (1 + t_{ij}) = \sum_G \prod_{\{i,j\} \in G} t_{ij}.$$

Here the base is all edges, the fiber is always  $\{0, 1\}$ . A section is a function from edges to  $\{0, 1\}$ , that is, the edges of a graph.

$$-1 \leq t_{ij} \leq 0.$$

## Example: The product rule

$$\prod_{i \in B} f_i(x+h) = \prod_{i \in B} (f_i(x) + f'_i(x)h + \epsilon_i)$$

$$\prod_{i \in B} f_i(x+h) = \sum_{R+S+T=B} \prod_{i \in R} f_i(x) \prod_{j \in S} f'_j(x)h \prod_{k \in T} \epsilon_k$$

$$\prod_{i \in B} f_i(x+h) = \prod_{i \in B} f_i(x) + \sum_{|S|=1} \prod_{i \in B \setminus S} f_i(x) \prod_{j \in S} f'_j(x)h + \epsilon$$

The derivative is

$$\sum_{|S|=1} \prod_{i \in B \setminus S} f_i(x) \prod_{j \in S} f'_j(x) = \sum_j \prod_{i \neq j} f_i(x) f'_j(x)$$

## Example: Fubini's theorem

$$\int f_1(s) ds \cdots \int f_n(x) ds = \int \cdots \int f(s_1) \cdots f(s_n) ds_1 \cdots ds_n.$$

$$\prod_i \int f_i(s) ds = \int \cdots \int \prod_i f(s_i) ds_1 \cdots ds_n.$$

# The connected graph sum

$$-1 \leq t_{ij} \leq 0$$

Goal: Estimate the connected graph sum

$$\sum_{G_c} \prod_{\{i,j\} \in E(G_c)} t_{ij}.$$

Estimate by tree sum:

$$\sum_T \prod_{\{i,j\} \in E(T)} |t_{ij}|.$$

Can we do better?



# Rooted trees as functions

- ▶ vertex set  $U$
- ▶ root  $r \in U$ .
- ▶  $\tau : U \setminus \{r\} \rightarrow U$

The orbit of each point in  $U$  ends in  $r$ .

# Connected graph identity

Take  $-1 \leq t_{ij} \leq 0$ .

This connected graph identity is a variant of an identity of Abdesselam and Rivasseau. Here  $\mathbf{s}$  is a family of integration variables in  $[0, 1]$ .

$$\sum_{G_c} \prod_{\{i,j\} \in G_c} t_{ij} = \sum_{\tau} c(\tau) W(\tau) \int I(\tau, \mathbf{s}) d\mathbf{s}.$$

$$c(\tau) = \prod_{i \neq r} t_{i\tau(i)}.$$

$$0 \leq W(\tau) \leq 1$$

$$0 \leq I(\tau, \mathbf{s}) \leq 1.$$

# Tree weights

- ▶ Withinlayer factor:

$$W(\tau) = \prod_{\{i,j\} | i \sim j \text{ mod } \tau} (1 + t_{ij}),$$

where the product is over two-element subsets  $\{i, j\}$ , and  $i \sim j \text{ mod } \tau$  means that  $i$  and  $j$  are at the same distance from the root in the tree  $\tau$ .

- ▶ Interlayer factor (random):

$$I(\tau, \mathbf{s}) = \prod_{i \neq r} \prod_{j \leftarrow i \text{ mod } \tau, j \neq \tau(i)} (1 + s_i t_{ij}),$$

where  $j \leftarrow i \text{ mod } \tau$  means that  $j$  is one tree distance unit closer to the root than the tree distance of  $i$  from the root.

For each vertex  $i \neq r$  there is a uniform integration variable  $s_i$  in the interval  $[0, 1]$  associated with the tree edge  $\{i, \tau(i)\}$ .

## Fix distances from the root

Decompose  $U$  into disjoint non-empty subsets.

$$\Delta = (U_0, U_1, \dots, U_h)$$

$$U_0 = \{r\}.$$

Let  $\mathcal{G}_c(\Delta)$  consist of all connected graphs  $G_c$  so the set of points in  $G_c$  a graph distance  $m$  from the root is  $U_m$ .

$$\sum_{G_c} \prod_{\{i,j\}} t_{ij} = \sum_{\Delta} \sum_{G_c \in \mathcal{G}_c(\Delta)} \prod_{\{i,j\} \in G_c} t_{ij}.$$

# 1 Distributive law: Resummation

$$\sum_{G_c \in \mathcal{G}_c(\Delta)} \prod_{\{i,j\} \in G_c} t_{ij} = W_\Delta l_\Delta,$$

Withinlayer factor:

$$W_\Delta = \prod_{i \sim j} (1 + t_{ij})$$

$i \sim j$  means  $i \neq j$  and  $i, j$  are at the same graph distance from the root.

Interlayer factor:

$$l_\Delta = \prod_{i \neq r} [\prod_{j \leftarrow i} (1 + t_{ij}) - 1].$$

$k \leftarrow i$  means the graph distance of  $k$  from the root is one less than the graph distance of  $i$  from the root.

## 2 Fundamental theorem of calculus: Product rule

Interlayer factor:

$$l_{\Delta} = \prod_{i \neq r} [\prod_{j \leftarrow i} (1 + t_{ij}) - 1].$$

$k \leftarrow i$  means the graph distance of  $k$  from the root is one less than the graph distance of  $i$  from the root.

FTC for each  $i \neq r$ :

$$\prod_{j \leftarrow i} (1 + t_{ij}) - 1 = \int_0^1 \frac{d}{ds} \prod_{j \leftarrow i} (1 + st_{ij}) ds = \sum_{k \leftarrow i} \int_0^1 t_{ik} \prod_{j \leftarrow i, j \neq k} (1 + st_{ij}) ds.$$

### 3 Distributive law: Trees emerge

The distributive law gives

$$\begin{aligned} I_{\Delta} &= \prod_{i \neq r} \sum_{k \leftarrow i} \left[ \int_0^1 t_{ik} \prod_{j \leftarrow i, j \neq k} (1 + st_{ij}) ds \right] \\ &= \sum_{\tau} \prod_{i \neq r} \left[ \int_0^1 t_{i\tau(i)} \prod_{j \leftarrow i, j \neq \tau(i)} (1 + st_{ij}) ds \right]. \end{aligned}$$

$\tau : U \setminus \{r\} \rightarrow U$

For  $i \neq r$  the value  $\tau(i) \leftarrow i$ .

## 4 Fubini's theorem

Let  $\mathcal{T}(\Delta)$  be the tree functions that respect decomposition  $\Delta$ .

$$I_{\Delta} = \sum_{\tau \in \mathcal{T}(\Delta)} \prod_{i \neq r} \left[ \int_0^1 t_{i\tau(i)} \prod_{j \leftarrow i, j \neq \tau(i)} (1 + st_{ij}) ds \right].$$

$$I_{\Delta} = \sum_{\tau \in \mathcal{T}(\Delta)} \int_0^1 \cdots \int_0^1 \prod_{i \neq r} \left[ t_{i\tau(i)} \prod_{j \leftarrow i, j \neq \tau(i)} (1 + s_i t_{ij}) ds_i \right].$$

Equivalently

$$I_{\Delta} = \sum_{\tau \in \mathcal{T}(\Delta)} c(\tau) \int I(\tau, \mathbf{s}) d\mathbf{s}.$$

$$c(\tau) = \prod_{i \neq r} t_{i\tau(i)}.$$

$$I(\tau, \mathbf{s}) = \prod_{i \neq r} \prod_{j \leftarrow i \text{ mod } \tau, j \neq \tau(i)} (1 + s_i t_{ij}),$$



# Conclusion

$$\sum_{G_c \in \mathcal{G}_c(\Delta)} \prod_{\{i,j\} \in G_c} t_{ij} = \sum_{\tau \in \mathcal{T}(\Delta)} c(\tau) W(\tau) \int I(\tau, \mathbf{s}) d\mathbf{s}.$$

Sum over  $\Delta$ .

$$\sum_{G_c} \prod_{\{i,j\} \in G_c} t_{ij} = \sum_{\tau} c(\tau) W(\tau) \int I(\tau, \mathbf{s}) d\mathbf{s}.$$

# Tree bound

Suppose that  $-1 \leq t_{ij} \leq 0$ . Then

$$\left| \sum_{G_c} \prod_{\{i,j\} \in G_c} t_{ij} \right| \leq \sum_{\tau} c^*(\tau).$$

Here

$$c^*(\tau) = \prod_{i \neq r} |t_{i\tau(i)}|.$$

Recursive structure

$$c^*(\tau_r) = \prod_{\tau(j)=r} |t_{rj}| c^*(\tau_j).$$

# Enriched tree bound

An approach to Fernández-Procacci bounds.

Suppose that  $-1 \leq t_{ij} \leq 0$ . Then

$$\left| \sum_{G_c} \prod_{\{i,j\} \in G_c} t_{ij} \right| \leq \sum_{\tau} C^*(\tau) = \sum_{\tau} c^*(\tau) Q(\tau). \quad (1)$$

The fiber weight is

$$Q(\tau) = \prod_{\tau(i)=\tau(j)} (1 + t_{ij}). \quad (2)$$

Recursive structure

$$C^*(\tau_r) = \left[ \prod_{\tau(j)=r} |t_{rj}| \prod_{\tau(i)=\tau(j)=r} (1 + t_{ij}) \right] \prod_{\tau(j)=r} C^*(\tau_j).$$

# Conclusion for physics

- ▶  $-1 \leq t_{ij} \leq 0$
- ▶  $|t_{ij}|$  measures the strength of repulsive interaction.
- ▶ The tree bound works when many interactions are weak, which happens at long distances.
- ▶  $0 \leq 1 + t_{ij} \leq 1$
- ▶ The enriched tree bound has additional factors  $1 + t_{ij}$ . is 1 when the interactions are large, which happens at short distances.
- ▶ As Fernández-Procacci show, such improved bounds give a better grasp of the physics.