

Correlation testing for affine invariant properties on F_p^n

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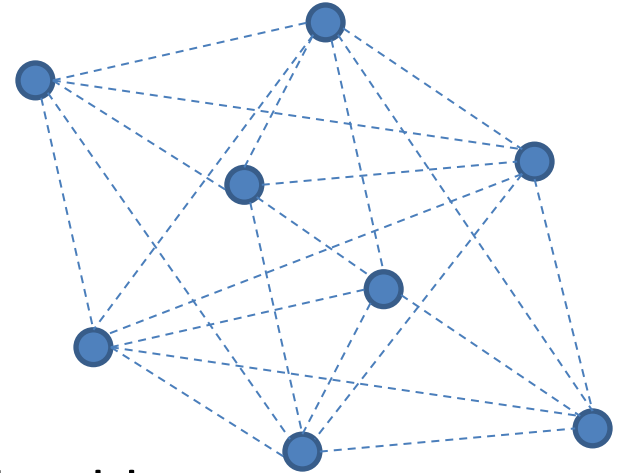
Joint with Hamed Hatami (McGill)

Property testing

- Math: infer **global structure** from **local samples**
- CS: Super-fast (randomized) algorithms for approximate decision problems
- Decide if large object **approximately has property**, while testing only a **tiny fraction** of it

Graph properties: 3-colorability

- Input: graph G
- Is G 3-colorable?
- Local test:
 - Sample $(1/\epsilon)^{O(1)}$ vertices
 - Accept if induced subgraph is 3-colorable
- Analysis:
 - Test always accepts 3-colorable graphs
 - Test rejects (w.h.p) graphs ϵ -far from 3-colorable



[Goldreich-Goldwasser-Ron'96]

Algebraic properties: linearity

- Input: function $f : \mathbb{F}_p^n \rightarrow \mathbb{F}_p$
- Is f **linear**?
- Local test:
 - Sample $x, y \in \mathbb{F}_p^n$
 - Check if $f(x + y) = f(x) + f(y)$
 - Repeat $1/\epsilon^{O(1)}$ times
- Analysis:
 - Test always **accepts linear functions**
 - Test **rejects** (w.h.p) functions **ϵ -far from linear**

0	1	3	0	2	5	1	2
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[Blum-Luby-Rubinfeld'90]

Codes: locally testable codes

- Code: $C \subset \mathbb{F}_p^n$
distinct elements have large distance
- Input: word $w \in \mathbb{F}_p^n$
- C is **locally testable** if there exists a (randomized) test which queries a few coordinates and
 - Always **accepts codewords**
 - **Rejects** (w.h.p) if w is **far from all codewords**
- The “mathematical core” of the PCP theorem
- Open: can C have constant rate, distance and testability?

Proofs: Probabilistic Checkable Proofs

- PCP Theorem: robust proof system
- Encoding of theorems + randomized local test (queries few bits of proof)
 - Test always **accepts legal proofs of theorems**
 - Test **rejects** (w.h.p) **proofs of false theorems**
- Major tool to prove hardness of approximation

Property testing: general framework

- Universe: set of objects (e.g. graphs)
- **Property**: subset of objects (e.g. 3-col graphs)
- Test: **randomized small sample** (e.g. small subgraph)

- Property is testable if **local consistency** implies **approximate global structure**

Which properties are testable?

- **Graph properties:** well understood
- **Algebraic properties:** partially understood
- **Locally testable codes:** major open problems
- **PCP / hardness of approximation:** whole field

Correlation testing

- Property testing: **strong** global structure
 - Is the object close to having property
- Correlation testing: **weak** global structure
 - Is the object slightly related to having property
- Motivation: possible generalizations of the inverse Gowers ~~conjecture~~ theorem

Correlation testing

Linearity correlation testing

- Function $f : \mathbb{F}_p^n \rightarrow \mathbb{F}_p$
- Correlation of f, g : $\langle f, g \rangle = \mathbb{E}_{x \in \mathbb{F}_p^n} [f(x) \overline{g(x)}]$
- Correlation with linear functions (characters):

$$\|\hat{f}\|_\infty = \max_{\ell: \mathbb{F}_p^n \rightarrow \mathbb{F}_p \text{ linear}} |\langle f, \omega_p^\ell \rangle|$$

$$(\omega_p = e^{2\pi i/p})$$

Linearity correlation testing

- Linear correlation: **global property**

Witnessed by **local average**

$$\begin{aligned} & \mathbb{E}_{x,y,z \in \mathbb{F}_p^n} [f(x+y+z) \overline{f(x+y)} \overline{f(x+z)} f(x)] \\ &= \sum_{\alpha} |\hat{f}(\alpha)|^4 \approx \|f\|_{U^2}^4 \end{aligned}$$

- Identifies functions correlated with linear funcs:
 - f correlated to linear: $|\hat{f}|_{\infty} \geq \varepsilon \implies \|f\|_{U^2} \geq \varepsilon$
 - f is not correlated: $|\hat{f}|_{\infty} \leq \delta \implies \|f\|_{U^2} \leq \sqrt{\delta}$

Linearity correlation testing

- Discrete setting: $f : \mathbb{F}_p^n \rightarrow \mathbb{F}_p$
Test queries **4 locations**, accepts f if
$$f(x + y + z) - f(x + y) - f(x + z) + f(x) = 0$$
- Acceptance probability:
 - **ϵ -correlated** with linear: prob. $\geq 1/p + \epsilon^2$
 - **negligible** correlation: prob. $\leq 1/p + o(1)$
- Property testing: **#queries** depends on ϵ
- Here: **#queries=4**, **acceptance prob.** depends on ϵ

Testing correlation with polynomials

- Inverse Gowers Theorem (for finite fields):

Global structure: correlation with low-degree polynomials (Higher-order Fourier coefs)

Witnessed by local average

Testing correlation with polynomials

- Correlation with **degree d polynomials**:

$$|f|_{u(Poly_d)} = \max_{Q: \mathbb{F}_p^n \rightarrow \mathbb{F}_p \text{ polynomial degree } \leq d} |\langle f, \omega_p^Q \rangle|$$

- Gowers norm: **average over 2^{d+1} points**

$$|f|_{U^{d+1}}^{2^{d+1}} = \mathbf{E}_{x, y_1, \dots, y_{d+1} \in \mathbb{F}_p^n} \left[\prod_{I \subseteq [d+1]} C^{d-|I|} f\left(x + \sum_{i \in I} y_i\right)\right]$$

$C = \text{Conjugation}$

Testing correlation with polynomials

- Direct theorem [Gowers]

$$\|f\|_{u(\text{Poly}_d)} \geq \varepsilon \implies \|f\|_{U^{d+1}} \geq \varepsilon$$

- Inverse Theorem [Bergelson-Tao-Ziegler]

$$\|f\|_{U^{d+1}} \geq \varepsilon \implies \|f\|_{u(\text{Poly}_d)} \geq \delta(\varepsilon)$$

(if $p < d$ then Poly_d = non classical polynomials)

Main theorem

- Gowers norms: **local averages** which witness global correlation to **low-degree polynomials**
- Question: are there other such **properties**?
 - Correlation witnessed by **local averages**
- Theorem [today]: **no**
(affine invariant properties, in large fields)

Correlation with property

- Property $P \subset \{g : \mathbb{F}_p^n \rightarrow \mathbb{F}_p\}$
(can also consider $P \subset \{g : \mathbb{F}_p^n \rightarrow \square\}$)
- Function $f : \mathbb{F}_p^n \rightarrow \mathbb{F}_p$
- Correlation of f with property P :

$$|f|_{u(P)} = \max_{g \in P} |\langle \omega^f, \omega^g \rangle|$$

Local test

- Local test (with q queries):
 - Distribution over $\{x_1, \dots, x_q\} \subset \mathbb{F}_p^n$
 - Local test $T : \mathbb{F}_p^q \rightarrow \{0, 1\}$
$$T(f) = \mathbb{E}[T(f(x_1), \dots, f(x_q))]$$
- T tests correlation with property P if
 - $\forall \varepsilon \exists \delta \in (0, \varepsilon), \theta^- < \theta^+$ such that
 - | $\|f\|_{u(P)} \geq \varepsilon \Rightarrow T(f) \geq \theta^+$
 - | $\|f\|_{u(P)} \leq \delta(\varepsilon) \Rightarrow T(f) \leq \theta^-$

Affine invariant properties

- Property $P \subset \{f : \mathbb{F}_p^n \rightarrow \mathbb{F}_p\}$

- P is affine invariant if

$$f(x) \in P \Leftrightarrow g(x) = f(Ax + b) \in P$$

- Examples:

- Linear functions; degree-d polynomials

- Functions with sparse / low-dim. Fourier representation

- **Local tests** for affine invariant properties are w.l.o.g **local averages over linear forms**

Local average over linear forms

- Variables $X = (X_1, \dots, X_k) \in (\mathbb{F}_p^n)^k$
- Linear form $L(X) = \lambda_1 X_1 + \dots + \lambda_k X_k \quad (\lambda_i \in \mathbb{F}_p)$
- System of linear forms $L = \{L_1, \dots, L_q\}$
 - E.g. $L = \{X + Y + Z, X + Y, X + Z, X\}$
- Average over linear forms:

$$T_{L, \alpha}(f) = \mathbf{E}_{X \in (\mathbb{F}_p^n)^k} \left[\omega_p^{\alpha_1 f(L_1(X)) + \dots + \alpha_q f(L_q(X))} \right]$$

$(\alpha \in \mathbb{F}_p^q)$

Local tests: affine invariant properties

- **Local tests** for affine invariant properties are w.l.o.g **averages over homogenous linear forms**
 - $L = \{L_1, \dots, L_q\}$ **homogenous** if $L_i(X) = X_1 + \sum_{i=2}^k \lambda_i X_i$
- \exists **systems of linear forms** L_i, α_i such that the sets

$$\{(T_{L_1, \alpha_1}(f), \dots, T_{L_m, \alpha_m}(f)) \parallel f \parallel_{u(P)} \geq \varepsilon\}$$

$$\{(T_{L_1, \alpha_1}(f), \dots, T_{L_m, \alpha_m}(f)) \parallel f \parallel_{u(P)} \leq \delta(\varepsilon)\}$$

are disjoint

Local tests: affine invariant properties

- Claim: any **local test** \rightarrow **local averages**

- Proof: P **affine invariant**, so $\forall A, b$

$$\|f\|_{u(P)} = \|f(Ax + b)\|_{u(P)}$$

- Choosing A, b **uniformly**:
 - transform each query (x_1, \dots, x_q)
 - to a **homogeneous system** $(Ax_1 + b, \dots, Ax_q + b)$

Main theorem (1)

- Property $P = (P_n \subset \{g : \mathbb{F}_p^n \rightarrow \mathbb{F}_p\})_{n \in \mathbb{N}}$
 - Consistent $P_n \subset P_{n+1}$
 - Affine invariant
 - Sparse $|P_n| = p^{o(p^n)}$
- Thm: If P is **locally testable** with q queries ($p > q$) then $\exists d \leq q$ such that for any sequence of functions

$(f_n : \mathbb{F}_p^n \rightarrow \mathbb{F}_p)_{n \in \mathbb{N}}$ which are unbiased $\lim_{n \rightarrow \infty} \mathbb{E} \omega_p^{f_n} = 0$

$$\liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0 \Leftrightarrow \liminf_{n \rightarrow \infty} \|f_n\|_{U^d} = 0$$

Main theorem (2)

- Consistent property

$$P = (P_n \subset \{g : \mathbb{F}_p^n \rightarrow \square, \|g\|_\infty \leq 1\})_{n \in \square}$$

- Thm: If P is **testable** by **systems of q linear forms** ($p > q$) then $\exists d \leq q$, for any bounded functions $(f_n : \mathbb{F}_p^n \rightarrow \square)_{n \in \square}$

$$\liminf_{n \rightarrow \infty} \|f_n - \mathbb{E}f_n\|_{u(P)} = 0 \Leftrightarrow \liminf_{n \rightarrow \infty} \|f_n - \mathbb{E}f_n\|_{U^d} = 0$$

- Q: Is this true for any **norm defined by linear forms**?

Proof

Main theorem

- $P = (P_n \subset \{g : \mathbb{F}_p^n \rightarrow \mathbb{F}_q, \|g\|_\infty \leq 1\})_{n \in \mathbb{N}}$
- P testable by systems of q linear forms ($q < p$)
- Thm: $u(P)$ norm equivalent to some U^d norm:
if $\lim_{n \rightarrow \infty} \mathbb{E} f_n = 0$ then

$$\liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0 \iff \liminf_{n \rightarrow \infty} \|f_n\|_{U^d} = 0$$

Proof idea

- Dfn: $S = \{\text{degrees } d: \forall \text{ large } n \exists \text{ degree-}d \text{ poly } Q_n$
 1. Q_n correlated with property P
 2. Q_n has “high enough” rank}
- $D = \text{Max}(S)$
 - D is bounded (bound depends on the linear systems)
- Lemma 1: $\liminf_{n \rightarrow \infty} \|f_n\|_{U^{D+1}} = 0 \implies \liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0$
- Lemma 2: $\liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0 \implies \liminf_{n \rightarrow \infty} \|f_n\|_{U^{D+1}} = 0$

Polynomial rank

- Q – degree d polynomial
- $\text{Rank}(Q)$ – minimal number of lower-degree polynomials R_1, \dots, R_c needed to compute Q
 - $Q(x) = \Gamma(R_1(x), \dots, R_c(x))$
- Thm [Green-Tao, Kaufman-L.]
If P has **high enough rank**, it has **negligible correlation with lower degree polynomials**

Polynomial factors

- Polynomial factor: $B = \{Q_1, \dots, Q_C : \mathbb{F}_p^n \rightarrow \mathbb{F}_p\}$
 - Sigma-algebra defined by Q_1, \dots, Q_C
 - $f : \mathbb{F}_p^n \rightarrow \square$: average over B, $E[f | B]$
- Complexity(B): C = number of basis polys
- Degree(B): max degree of Q_1, \dots, Q_C
- Rank(B): min. rank of linear comb. of Q_1, \dots, Q_C
 - Large rank: $Q_1(x), \dots, Q_C(x)$ are **nearly independent**

Decomposition theorems

- Fix $d < p$
- $f : \mathbb{F}_p^n \rightarrow \mathbb{R}$ can be decomposed as $f = f_1 + f_2$
 - $f_1 = \mathbb{E}[f \mid B]$
B has degree d , high rank, bounded complexity
 - $\|f_2\|_{U^{d+1}} \leq \epsilon$

Complexity of linear systems

- Linear form: $L(X) = \lambda_1 X_1 + \dots + \lambda_k X_k$
- Linear system: $\mathbf{L} = \{L_1, \dots, L_q\}$
- **Average:** $T_{\mathbf{L}}(f) = \mathbb{E}_{\mathbf{X} \in (\mathbb{F}_p^n)^k} \prod_{i=1}^q f(L_i(X))$
- **Complexity:** min. d , if $f = f_1 + f_2$, $\|f_2\|_{U^{d+1}} \leq 1$
then $T_{\mathbf{L}}(f) \approx T_{\mathbf{L}}(f_1)$
- C-S complexity [Green-Tao]
- True complexity [Gowers-Wolf, Hatami-L.]

Proof idea

- Dfn: $S = \{\text{degrees } d: \forall \text{ large } n \exists \text{ degree-}d \text{ poly } Q_n$
 1. Q_n correlated with property P
 2. Q_n has “high enough” rank}
- $D = \text{Max}(S)$
 - D is bounded (\leq complexity of linear systems)
- Lemma 1: $\liminf_{n \rightarrow \infty} \|f_n\|_{U^{D+1}} = 0 \implies \liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0$
- Lemma 2: $\liminf_{n \rightarrow \infty} \|f_n\|_{u(P)} = 0 \implies \liminf_{n \rightarrow \infty} \|f_n\|_{U^{D+1}} = 0$

Lemma 1: Small $U^{D+1} \rightarrow$ small $u(P)$

- D : max deg of **high rank polys correlate with P**
- Assume $\|f\|_{U^{D+1}} \approx 1$ but $\|f\|_{u(P)} \geq \varepsilon$

- Step 1: reduce to “**structured function**”

- Linear system L of complexity S ($S > D$)

- Decompose: $f = f_1 + f_2$

$$f_1 = E[f | B], \|f_2\|_{U^{S+1}} \approx 1$$

- Reduce to studying f_1 - func. of **deg $\leq S$ polys**:

- $\|f_1\|_{U^{D+1}} \approx 1$

- $T_L(f) \approx T_L(f_1) \Rightarrow \|f_1\|_{u(P)} \geq \varepsilon'$

Lemma 1: Small $U^{D+1} \rightarrow$ small $u(P)$

- D : max deg of high rank polys correlate with P
- Structured function: $f_1 = E[f | B]$, $\deg(B) \leq S$
 - $\|f_1\|_{U^{D+1}} \leq 1$
 - $\|f_1\|_{u(P)} \geq \varepsilon'$
- Will show: $\|f_1\|_{u(P)} \approx 0$
- Use the structure: $f_1(x) = \sum \alpha_i \omega_p^{Q_i(x)}$, $\deg Q_i \leq S$
 - $\deg(Q_i) \leq D \Rightarrow \alpha_i \approx 0$ because $\|f_1\|_{U^{D+1}} \leq 1$
 - $\deg(Q_i) > D \Rightarrow \|\omega_p^{Q_i}\|_{u(P)} \approx 0$ by def of D

Lemma 2: small $u(P) \rightarrow$ small U^{D+1}

- Key ingredient: **invariance principle**
 - High rank polynomials “look the same” to averages

- $\{Q_1, \dots, Q_c\}, \{Q'_1, \dots, Q'_c\}$ high rank, $\deg(Q_i) = \deg(Q'_i)$

$$f(x) = \Gamma(Q_1(x), \dots, Q_c(x))$$

$$f'(x) = \Gamma(Q'_1(x), \dots, Q'_c(x))$$

Then local averages **cannot distinguish f, f'** :

$$T_L(f) \approx T_L(f')$$

Part 2: small $u(P) \Rightarrow$ small U^{D+1}

- D : max deg of high rank polys correlate with P
- Assume $\|f\|_{u(P)} \approx 1$ but $\|f\|_{U^{D+1}} \geq \varepsilon$
 - Reduce to structured function, $f_1 = E[f | B]$
- f_1 correlated with high-rank Q of degree $\leq D$
 - Assume for now: $\deg(Q)=D$
- Dfn of D : Exists high rank poly Q' , $\deg(Q')=D$,
 Q' correlated with some function $g \in P$
- Contradiction: Define $f'_1 = f_1$ with Q replaced by Q'
 - Invariance principle: $T_L(f_1) \approx T_L(f'_1)$
 - f'_1 is correlated with $g \in P$

Part 2: small $u(P) \rightarrow$ small U^{D+1}

- Problem: what if f_1' correlated with high rank poly of **degree $< D$** ?
 - Solution: can find Q' correlated with property P for of **all degrees $\leq D$**
 - Reason: **systems of averages** are **robust**
- Thm: for any **family of linear systems**, the set
$$\{(T_{L_1}(f), \dots, T_{L_k}(f)) : f : \mathbb{F}_p^n \rightarrow \mathbb{F} \mid \|f\|_\infty \leq 1\} \subset \mathbb{F}^k$$
has a **non-empty interior** for some finite n (unless not for trivial reasons)
 - analog of [Erdos-Lovasz-Spencer] for additive settings

Summary

- Property testing: witness **strong structure** by **local samples**
- Correlation test: witness **weak structure**
- Main result: any **affine invariant property** which is correlation testable, is **essentially equivalent** to **low-degree polynomials**

Open problems

- Which **norms** can be defined by local averages
 - Are always equivalent to some U^d norm?
- Testing in **low characteristics**
- Is it possible to test if a function $f : \mathbb{F}_2^n \rightarrow \mathbb{F}_2$ is correlated with cubic polynomials?
 - U^4 norm doesn't work
 - Unknown even if #queries depends on correlation

THANK YOU!