

# Orthogonal Latin Hypercube Designs for Computer Experiments

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This talk is based on two papers:

- Lin, Mukerjee and Tang (2009 *Biometrika*)
- Lin, Bingham, Sitter and Tang (2010 *Annals of Statistics*)

# Outline

- 1 Introduction
- 2 Lin, Mukerjee and Tang (2009)
- 3 Lin, Bingham, Sitter and Tang (2010)
- 4 Concluding remarks

# Introduction

## Designs for Computer Experiments

- 1 Model-dependent criteria (Santner, Williams and Notz 2003)
- 2 Space-filling designs
  - Maximin distance designs  
(Johnson, Moore and Ylvisaker 1990)
  - OA based designs (Owen 1992; Tang 1993)
  - $(t, m, s)$ -nets from quasi Monte Carlo (Niederreiter 1992)
- 3 **Orthogonal Latin hypercubes** (Ye 1998; Butler 2001)

# Introduction

## Orthogonal Latin Hypercube Designs

- 1 Why Latin hypercube designs?
  - guarantee one-dimensional stratification
  - a simple structure for finding good designs
- 2 Why orthogonality?
  - directly useful for regression type modeling
  - provides a stepping stone to space-filling

# Introduction

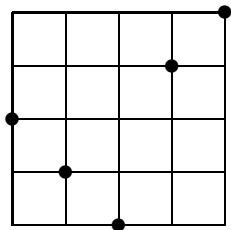
## Latin Hypercube Designs

- 1 An  $n \times m$  matrix  $D = (d_{ij})$  is called a *Latin hypercube* if each column of  $D$  is a permutation of  $1, \dots, n$ .
- 2 Two ways of generating design points within  $[0, 1]^m$ 
  - $x_{ij} = (d_{ij} - 0.5)/n$ ,                       $x_{ij} = (d_{ij} - u_{ij})/n$
  - the  $n$  points given by  $(x_{i1}, \dots, x_{im})$  with  $i = 1, \dots, n$

## Introduction

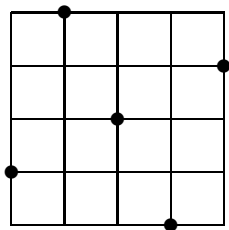
$$D_1$$

1	3
2	2
3	1
4	4
5	5



$$D_2$$

1	2
2	5
3	3
4	1
5	4



# Lin, Mukerjee and Tang (2009)

- Use the idea of rotation. Early work includes Steinberg and Lin (2006) and Pang, Liu and Lin (2009).
- Let  $B = (b_{ij})$  be a Latin hypercube of  $n$  runs for  $p$  factors with each column being a permutation of  $-(n-1)/2, -(n-3)/2, \dots, (n-3)/2, (n-1)/2$ .
- Let  $A$  be an  $OA(n^2, 2f, n, 2)$ . Denote the symbols in  $A$  by  $1, 2, \dots, n$ .



# Lin, Mukerjee and Tang (2009)

The construction has three steps.

**Step 1.** For each  $j$ , obtain  $A_j$  from  $A$  by replacing  $1, 2, \dots, n$  by  $b_{1j}, b_{2j}, \dots, b_{nj}$  respectively, and then partition  $A_j$  as  $A_j = [A_{j1}, \dots, A_{jf}]$ , where each of  $A_{j1}, \dots, A_{jf}$  has two columns.

**Step 2.** For each  $j$ , obtain  $M_j = [A_{j1}V, \dots, A_{jf}V]$ , where

$$V = \begin{bmatrix} 1 & -n \\ n & 1 \end{bmatrix}.$$

**Step 3.** Finally, obtain  $M = [M_1, \dots, M_p]$ , of order  $N \times q$ , where  $N = n^2$  and  $q = 2pf$ .

# Lin, Mukerjee and Tang (2009)

If  $p = 1$ , the above construction is equivalent to that in Steinberg & Lin (2006) and Pang et al. (2009). The power of our method is derived from allowing  $p \geq 2$ .

**Theorem 1.** *We have that*

- *$M$  is a Latin hypercube;*
- *$M$  is orthogonal if  $B$  is orthogonal.*
- *$M$  is nearly orthogonal if  $B$  is nearly orthogonal.*

# Lin, Mukerjee and Tang (2009)

## Some example

We use  $OLH(n, p)$  to denote an orthogonal Latin hypercube with  $n$  rows and  $p$  columns.

The orthogonal Latin hypercubes shown in Table 1 were obtained by a computer search. To the best of our knowledge, these are new.

# Lin, Mukerjee and Tang (2009)

Table 1. *Small Orthogonal Latin hypercubes*

OLH(7, 3)	OLH(9, 5)	OLH(11, 7)										
-3	3	2	-4	-2	0	-3	3	-5	-4	-5	-5	-
-2	0	-3	-3	4	2	1	-2	-4	2	-1	3	-
-1	-2	-1	-2	-3	-4	-1	-3	-3	-2	4	5	-
0	-3	1	-1	3	-2	3	4	-2	3	-3	4	-
1	-1	3	0	-4	4	4	0	-1	4	2	-4	-
2	1	-2	1	2	-1	0	-4	0	-5	5	-2	-
3	2	0	2	0	3	-2	-1	1	5	3	-3	-
			3	1	1	-4	2	2	-1	1	1	-
			4	-1	-3	2	1	3	0	0	-1	-
								4	1	-4	0	-
								5	-3	-2	2	-

In addition to the above, we will also use the  $OLH(5, 2)$  and  $OLH(8, 4)$  constructed by Ye (1998).

# Lin, Mukerjee and Tang (2009)

- Let  $B$  be  $OLH(5, 2)$ ,  $OLH(7, 3)$ ,  $OLH(8, 4)$ ,  $OLH(9, 5)$  and  $OLH(11, 7)$
- Let  $A$  be  $OA(25, 6, 5, 2)$ ,  $OA(49, 8, 7, 2)$ ,  $OA(64, 8, 8, 2)$ ,  $OA(81, 10, 9, 2)$  and  $OA(121, 12, 11, 2)$ .
- We obtain  $OLH(25, 12)$ ,  $OLH(49, 24)$ ,  $OLH(64, 32)$ ,  $OLH(81, 50)$ , and  $OLH(121, 84)$ .

# Lin, Bingham, Sitter and Tang (2010)

Consider the following construction:

$$L = A \otimes B + n_2 C \otimes D,$$

- $A = (a_{ij})_{n_1 \times m_1}$ , with  $a_{ij} = \pm 1$ ,
- $B_{n_2 \times m_2}$  is a Latin hypercube
- $C_{n_1 \times m_1}$  is a Latin hypercube,
- $D = (d_{ij})_{n_2 \times m_2}$ , with  $d_{ij} = \pm 1$ .

Design  $L$  has  $n = n_1 n_2$  runs and  $m = m_1 m_2$  factors.

# Lin, Bingham, Sitter and Tang (2010)

**Theorem 2.** *Design  $L$  is an orthogonal Latin hypercube if*

- (i)  $A$  and  $D$  are orthogonal matrices of  $\pm 1$ ,*
- (ii)  $B$  and  $C$  are orthogonal Latin hypercubes,*
- (iii)  $A^T C = 0$  or  $B^T D = 0$ ,*
- (iv) there does not exist  $i$  such that  $a_{pi} = -a_{p'i}$  where  $p$  and  $p'$  are determined by  $c_{pi} = -c_{p'i}$ .*

( $L$  is nearly orthogonal if  $B$  and  $C$  are nearly orthogonal)

# Lin, Bingham, Sitter and Tang (2010)

**Theorem 3.** Suppose that an  $\text{OLH}(n, m)$  is available, where  $n$  is a multiple of 4 such that a Hadamard matrix of order  $n$  exists. Then we can construct  $\text{OLH}(2n, m)$ ,  $\text{OLH}(4n, 2m)$ ,  $\text{OLH}(8n, 4m)$ ,  $\text{OLH}(16n, 8m)$ .



## Lin, Bingham, Sitter and Tang (2010)

Theorem 3 is a very powerful result. By repeated application, one can obtain many infinite series of orthogonal Latin hypercubes.

For example, starting with an  $OLH(12, 6)$ , we can obtain an  $OLH(192, 48)$ , which can be used in turn to construct an  $OLH(768, 96)$  and so on.

For another example, an  $OLH(256, 248)$  in Steinberg and Lin (2006) can be used to construct an  $OLH(1024, 496)$ , an  $OLH(4096, 1984)$  and so on.

# Lin, Bingham, Sitter and Tang (2010)

For any run size  $n$ , what is the maximum number  $m^*$  of factors for an  $\text{OLH}(n, m^*)$  to exist?

**Theorem 4.** *Let  $n = 16k + j$ . Then we have that:*

(i)  $m^* = 1$  if  $j = 2, 6, 10, 14$ ;

(ii)  $m^* \geq 6$  for all  $n$  where  $k \geq 1$  and  $j \neq 2, 6, 10, 14$ ;

(iii)  $m^* \geq 7$  for  $n = 16k + 11$  where  $k \geq 0$ ;

(iv)  $m^* \geq 12$  for  $n = 16k, 16k + 1$  where  $k \geq 2$ ;

(v)  $m^* \geq 24$  for  $n = 32k, 32k + 1$  where  $k \geq 2$ ;

(vi)  $m^* \geq 48$  for  $n = 64k, 64k + 1$  where  $k \geq 2$ .

# Lin, Bingham, Sitter and Tang (2010)

Note that

- ① Part (i) is a non-existence result.
- ② Other parts of the theorem make use of a stacking method in combination with Table 2.

*Table 2. The maximum number  $m$  of columns in  $OLH(n, m)$  by the algorithm for  $4 \leq n \leq 21$*

$n$	4	5	7	8	9	11	12	13	15	16	17	19	20	21
$m$	2	2	3	4	5	7	6	6	6	12	6	6	6	6

# Concluding remarks

- Discussed two general methods for constructing orthogonal or nearly orthogonal Latin hypercubes.
- Powerful methods both in terms of the large ratio  $m/n$  and the flexibility of the run size.
- Future work?